# Discrete Differential Geometry

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May 2013

Images courtesy of Konrad Polthier, Wikipedia

## Introduction

Discrete differential geometry

- Deals with discrete geometric objects: points (vertices), segments (edges), faces, polygons, polyhedra.
- Finds properties and invariants (affine or metric)
- Focuses on differential quantities (such as curvature)



How do we discretize derivatives?

Tangent vectors (s is arclength)

$$t(s) = \frac{d\gamma}{ds} = \lim_{h \to 0} \frac{\gamma(s+h) - \gamma(s)}{h} \longleftrightarrow t_i = \frac{p_{i+1} - p_i}{\ell_i}$$

• Curvature in  $\mathbb{R}^2$  (change in the tangent direction)

$$\kappa(s)n(s) = \frac{dt}{ds} \stackrel{?}{\longleftrightarrow} \kappa_i n_i = \frac{t_{i+1} - t_i}{\ell_i}$$

(note the lack of symmetry!)

• Curvature of surfaces in  $\mathbb{R}^3$ ?

## Smooth gaussian curvature

Gaussian curvature K at p measures

Length of infinitesimal circle at p:

$$K = \lim_{r \to 0} \frac{3}{\pi r^3} (2\pi r - L(C(r)))$$

variation of the normal vector w.r.t. the metric

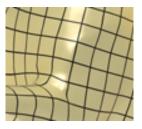
$$K = \frac{\det II}{\det I} = \frac{\begin{vmatrix} -\frac{E_{vv}}{2} + F_{uv} - \frac{G_{uu}}{2} & \frac{E_u}{2} & F_u - \frac{E_v}{2} \\ F_v - \frac{G_u}{2} & E & F \\ \frac{G_v}{2} & F & G \end{vmatrix} - \begin{vmatrix} 0 & \frac{E_v}{2} & \frac{G_u}{2} \\ \frac{E_v}{2} & E & F \\ \frac{G_u}{2} & F & G \end{vmatrix}}{(EG - F^2)^2}$$

spherical area of the normal

Curvature and derivatives in multiple variables

Problems:

- what are I and II?
- how can one take derivatives when there are no notion of "two variables"?



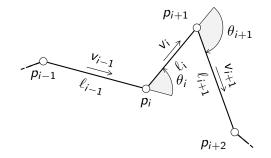


Conclusion: describe geometric invariants by their properties.

#### Discrete curves in the plane

A (connected) curve is set of vertices  $p_1, \ldots, p_n$  linked by edges  $(p_i p_{i+1})$ . Its (euclidean) invariants are:

- ▶ lengths  $\ell_i = \|p_{i+1} p_i\|$ , defined on edges
- tangent vectors  $v_i = \frac{p_{i+1}-p_i}{\ell_i}$
- rotation angles  $\theta_i = (v_{i-1}, v_i)$



## Bonnet theorem

#### Theorem (Bonnet theorem, smooth case)

Given a function  $\kappa$ , there exists a curve  $\gamma$  whose curvature is exactly  $\kappa$ . Such a curve is unique up to congruence.

Solve in arclength coordinate s, using Cauchy–Lipschitz theorem:

$$\frac{d^2\gamma}{ds^2} = \kappa J \frac{d\gamma}{ds},$$

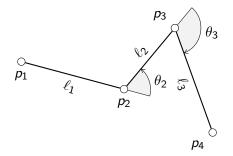
where J is the rotation by  $+\pi/2$ . The (local) solution is determined by initial data  $\gamma(0)$ ,  $\gamma'(0)$ , hence up to congruence (equivariance of the equation).

Global existence is more complicated.

# Bonnet theorem (discrete case)

#### Theorem

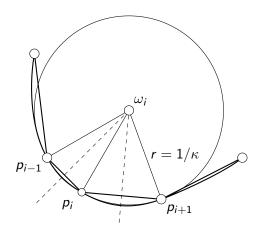
Given angles  $\theta_i$  and lengths  $\ell_i$ , there exists a discrete curve  $(p_i)$  corresponding to these. Such a curve is unique up to congruence.



#### Curvature as closest fitting circle

Menger curvature:

$$\kappa_i^M = \frac{2\sin\theta_i}{\|p_{i+1} - p_{i-1}\|}$$



#### Variation of tangent or normal vectors

In the smooth case,  $\kappa = \langle dt/ds, Jt \rangle = \frac{d\phi}{ds}$ , where  $\phi$  is the angle of the tangent t (or normal n). For discrete curves, since  $\phi_i - \phi_{i-1}$ ,

$$\kappa_i^{\theta} = \frac{2\theta_i}{\ell_{i-1} + \ell_i} \quad \text{or} \quad \kappa_i^n = \frac{2\langle t_i - t_{i-1}, n_i \rangle}{\ell_{i-1} + \ell_i} = \frac{4\sin\frac{\theta_i}{2}}{\ell_{i-1} + \ell_i}$$

As a consequence, we recover Gauss-Bonnet

$$\theta(b) - \theta(a) = \int_{a}^{b} \kappa(s) \, ds$$
  
total angular variation  $= \sum_{i} \theta_{i} = \sum_{i} \kappa_{i}^{\theta} \, \frac{\ell_{i-1} + \ell_{i}}{2}$ 

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## Variation of length

Consider the length  $L = \sum_{i} \ell_{i} = \sum_{i} ||p_{i} - p_{i-1}||$ .

$$abla_{p_i} L = 
abla_{p_i} \| p_i - p_{i-1} \| + \| p_{i+1} - p_i \| = v_{i-1} - v_i$$

Mean curvature vector  $\vec{H_i} := \frac{2(v_i - v_{i-1})}{\ell_{i-1} + \ell_i} = \kappa_i^L n_i$  with

$$\kappa_i^L = \kappa_i^n = \frac{4\sin\frac{\theta_i}{2}}{\ell_{i-1} + \ell_i}$$

First variation of area, when  $p_i 
ightarrow p_i' = p_i + \epsilon u_i$  is

$$L' = L - \epsilon \sum_{i} \kappa_{i}^{L} \langle u_{i}, n_{i} \rangle + o(\epsilon)$$

#### Approximation properties

Convergence when points are close enough?

• 
$$\gamma$$
 is the graph of  $x \mapsto y(x) = \frac{c}{2}x^2 + o(x^2)$   
•  $p_{i-1} = (-\tau_{i-1}, \frac{c}{2}\tau_{i-1}^2 + o(\tau_{i-1}^2)), p_{i+1} = (\tau_i, \frac{c}{2}\tau_i^2 + o(\tau_i^2))$  with  
positive  $\tau_i, \tau_{i-1} \in O(\tau), \tau \to 0$ .  
 $\ell_i = \tau_i \left(1 + \frac{c^2}{8}\tau_i^2 + o(\tau_i^2)\right), \quad \ell_{i-1} = \tau_{i-1} \left(1 + \frac{c^2}{8}\tau_{i-1}^2 + o(\tau_{i-1}^2)\right)$   
 $\overrightarrow{p_{i-1}p_i} \cdot \overrightarrow{p_ip_{i+1}} = \tau_{i-1}\tau_i \left(1 - \frac{c^2}{4}\tau_i\tau_{i-1} + o(\tau_{i-1}\tau_i)\right)$   
 $\cos \theta_i = \frac{\overrightarrow{p_{i-1}p_i} \cdot \overrightarrow{p_ip_{i+1}}}{\ell_i\ell_{i-1}} = \frac{\left(1 - \frac{c^2}{4}\tau_i\tau_{i-1} + o(\tau_{i-1}\tau_i)\right)}{\left(1 + \frac{c^2}{8}\tau_i^2 + o(\tau_i^2)\right) \left(1 + \frac{c^2}{8}\tau_i^2 + o(\tau_i^2)\right)}$   
 $= 1 - \frac{c^2}{8}(\tau_i + \tau_{i-1})^2 + o(\tau^2) = 1 - \frac{\theta_i^2}{2} + o(\tau^2)$ 

$$\theta_i \sim \frac{c}{2} \left( \tau_{i-1} + \tau_i \right)$$

### Discrete surfaces

An (abstract) discrete (polyhedral) surface S is

- a set V of vertices
- ▶ a set *E* of (unordered) edges e = (pq),  $p, q \in V$ ,  $p \neq q$

▶ a set *F* of faces  $f = (p_1, ..., p_n = p_1)$ , where  $(p_i p_{i+1}) \in E$  such that

- every vertex *p* belongs to 3 edges (2 on the boundary)
- every edge e belongs to 2 faces (1 on the boundary)
- every face f has at least 3 vertices (edges)
- Ineighbors of p form a closed curve (except on the boundary)
- boundary curves are closed

#### Metric structure

A metric structure on S is given by

- identification of each face f with a (non-unique) planar polygon (chart), with compatible edge length
- by the lengths ℓ(e) of each edge e, and angles between adjacent edges, denoted α<sub>f</sub>(p).

Remarks:

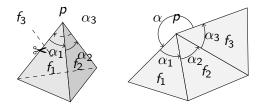
- for a triangle, lengths suffice
- for general polygons, lengths and angles are needed, and they have to satisfy a integrability relation

$$\sum_{j=1}^n \ell_j e^{i \sum_{k=1}^j \theta_k} = 0$$

where  $\theta_k = \pi - \beta_k$  is the rotation angle.

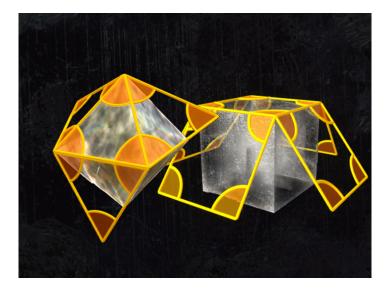
# Intrinsic distance and unfolding

Unfolding = constructing isometric faces with connected edges.



The intrinsic metric inside S is flat with conical singularities ( $\Rightarrow$  straight lines across edges). The conical singularity is measured by the angle defect (impossibility to close the unfolded faces)  $\alpha = 2\pi - \sum \alpha_k$ . Also recover gaussian curvature via : length of circle around p.

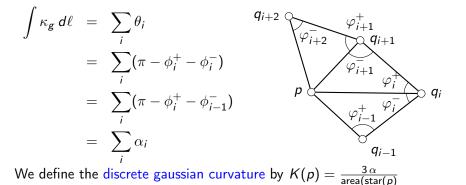
## Unfolding preserves intrinsic data



#### Gauss–Bonnet around *p*

Apply the Gauss–Bonnet formula to faces touching p (aka star(p)):

$$\int K \, dA + \int \kappa_{g} \, d\ell = 2\pi$$



Global discrete Gauss-Bonnet

We can directly prove

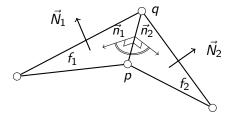
$$\sum_{p} K(p) = 2\pi \chi = 2\pi (|V| - |E| + |F|)$$

For a surface without boundary,

$$\sum_{p} \alpha(p) = \sum_{p} (2\pi - \sum_{f \ni p} \alpha_{f}(p)) = 2\pi |V| - \sum_{f} \sum_{p \in f} \alpha_{f}(p)$$
  
=  $2\pi |V| - \sum_{f} (|f| - 2)\pi = 2\pi |V| + 2\pi |F| - \pi \sum_{f} |f|$   
=  $2\pi (|V| - |E| + |F|)$ 

## Dihedral angles

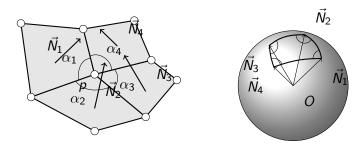
Gaussian curvature relies on intrinsic distances, but extrinsic position is determined by dihedral angles  $\theta = (\vec{N}_1, \vec{N}_2)$ .



#### Theorem (Reconstruction)

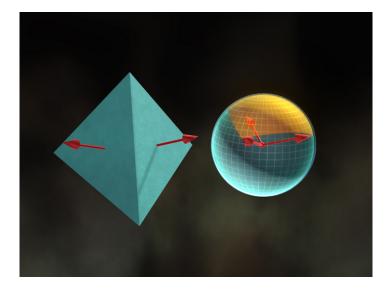
A surface in space is uniquely determined by its lengths, vertex angles and diedral angles, up to congruence.

## Spherical area



- spherical distance is equal to the dihedral angle
- spherical vertex angle at  $\vec{N}_i$  is  $\pi \alpha_i$

## Gauss image of the top vertex



### Variation of area

The area of triangle f = (pqr) is  $A(f) = \frac{1}{2} \overrightarrow{qp} \cdot J \overrightarrow{qr}$ . So, the partial gradient of area of a triangulated polyhedral surface S is

$$\nabla_{p}A = \frac{1}{2} \sum_{(pqr)\in F} J \overrightarrow{qr} = \frac{1}{2} \sum_{i} J_{i} \overrightarrow{q_{i}q_{i+1}}$$

Cotangent formula  

$$J\overline{q_{i}q_{i+1}} = \cot \varphi_{i+1}^{-} \overrightarrow{q_{i}\rho} + \cot \varphi_{i}^{+} \overrightarrow{q_{i+1}\rho}$$

$$\nabla_{p}A = \frac{1}{2} \sum_{i=1}^{n} (\cot \varphi_{i+1}^{-} + \cot \varphi_{i-1}^{+}) \overrightarrow{q_{i}\rho}$$

$$q_{i+2}$$

$$\varphi_{i+1}^{-} q_{i+1}$$

$$\varphi_{i+1}^{-} \varphi_{i+1}^{-} \varphi$$

#### Mean curvature vector

$$\vec{H}(p) = -\frac{1}{2} \nabla_p A = -\frac{1}{4} \sum_{e \sim p} J \vec{e} = -\frac{1}{4} \sum_i J_i \overrightarrow{q_i q_{i+1}}$$

$$= -\frac{1}{4} \sum_i J_i (\overrightarrow{q_i p} - \overrightarrow{q_{i+1} p}) = -\frac{1}{4} \sum_i (J_i - J_{i-1}) \overrightarrow{q_i p}$$

$$= \frac{1}{4} \left( \sum_i \vec{N_i} \times \overrightarrow{pq_i} + \sum_i \vec{N_{i-1}} \times \overrightarrow{q_i p} \right) = \frac{1}{4} \sum_i (\vec{N_{i-1}} - \vec{N_i}) \times \overrightarrow{pq_i}$$

$$= \frac{1}{2} \sum_i |e_i| \sin \frac{\theta_i}{2} \vec{N_{e_i}} = \frac{1}{2} \sum_{e \ni p} |e| \sin \frac{\theta_e}{2} \vec{N_e}$$

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Solve the smoothing problem by moving the points to minimize area (hence bumps):

$$\frac{dp_i}{dt} = \vec{H}(p_i)$$

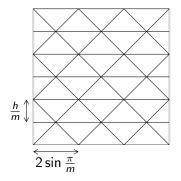
# Schwarz lantern (1890)

Discretize the cylinder of radius r and height h with integers n, m

- ► 2*nm* triangles
- distance to cylinder bounded by

$$r(1-\cos\frac{\pi}{n}) \xrightarrow[n\to\infty]{} 0$$

- area of a single triangle :  $\frac{rh}{m}\sin\frac{\pi}{n}\sqrt{1+\frac{4r^2}{h^2}m^2\sin^4\frac{2\pi}{n}}$
- ► total area  $2rh\left(n\sin\frac{\pi}{n}\right)\sqrt{1+\frac{4r^2}{h^2}m^2\sin^4\frac{2\pi}{n}}$
- ► limit when  $n \to \infty$ ,  $\frac{m}{n^2} \to \alpha \le \infty$ :  $2\pi rh \sqrt{1 + \frac{16\pi^4 r^2}{h^2} \alpha^2}$
- ► vertical slope of triangles of angle approx  $\frac{2h}{\alpha\pi^2 r}$



#### Discrete curvature

Xu, Xu & Sun (2005) prove that there is no discrete expression that will converge to curvature for all discretization Proof :

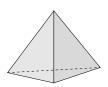
- ▶ fix points p = (0,0) ,  $q_i = (\epsilon x_i, \epsilon y_i, \epsilon z_i)$ ,  $i = 1, \dots, 4$
- ▶ consider a surface S given by  $z = \frac{1}{2}(ax^2 + 2bxy + cy^2)$ , then the gaussian curvature is  $K = ac b^2$  and obviously  $p \in S$
- ▶ solve  $q_i \in S$  in terms of a, b, c, i.e.  $ax_i^2 + 2bx_iy_i + c^2y_i^2 = 2z_i/\epsilon$ the rank may lower than 3, e.g.  $(x_1, y_1) = (1, 0) = -(x_3, y_3)$  and  $(x_2, y_2) = (0, 1) = -(x_4, y_4)$ ; then b is arbitrary.

$$\begin{pmatrix} x_1^2 & 2x_1y_1 & y_1^2 \\ x_2^2 & 2x_2y_2 & y_2^2 \\ x_1^3 & 2x_3y_3 & y_3^2 \\ x_4^2 & 2x_4y_4 & y_4^2 \end{pmatrix}$$

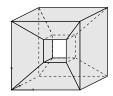
Why is the *discrete* Euler characteristic a topological invariant?

$$\chi = |V| - |E| + |F|$$

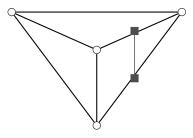
- add/remove vertex on edge, edge between vertices, vertex in face
- $\blacktriangleright\,\Rightarrow$  always assume triangulation if needed
- ▶ define an equivalence relation : S ~ S' if they can be refined up to the same combinatorial structure. Then \(\chi(S) = \chi(S')\).
- Is is trivial?

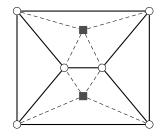


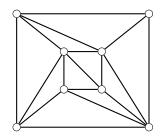




## From the tetrahedron to the cube







Pascal Romon (UPEMLV)

# Smooth Euler characteristic

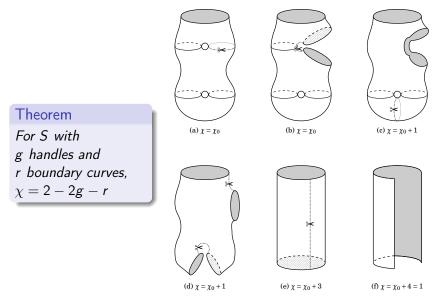
Define  $\chi$  on smooth S by picking a triangulation (or tiling). Why is it well-defined?

- invariant under homeomorphisms (bicontinuous bijections)
- move vertices and edges
- $\blacktriangleright$  refining to prove that  $\chi$  is independent of the tiling

Notes:

- for extrinsic surfaces, one can deform from smooth surfaces to discrete polyhedra
- $\blacktriangleright$  there are other meanings of  $\chi$  . . .

## Cuts and handles



## Total curvature

#### Theorem (discrete Gauss-Bonnet)

 $\sum lpha(\mathbf{p}) = 2\pi \chi$  (where  $lpha(\mathbf{p})$  is the angle defect)

Apply to the equilateral metric:

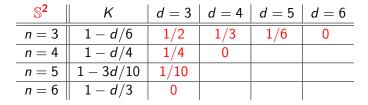
- ► faces are regular *n*-gons
- edge length = 1
- inside angle is  $(1-2/n)\pi$  (Gauss-Bonnet again)
- angle defect is  $\alpha(p) = 2\pi K_c(p)$  (combinatorial curvature)

$$\mathcal{K}_c(p) = 1 - \frac{d(p)}{2} + \sum_{f \ni p} \frac{1}{|f|}$$

# Regular tilings of $\mathbb{S}^2$ and $\mathbb{T}^2$

A tiling is *regular* if degree d and number n = |f| of edges are constant.

$$\chi = \sum_{p \in V} K_c(p) = |V| \left(1 - d \frac{n-2}{2n}\right)$$





#### Asymptotic degree

Consider tilings with constant number of edges n = |f|. Let  $v_k$  be the number of vertices of degree k. Then

$$|V| = \sum_{k} v_{k}, 2|E| = n|F|$$

$$|E| = \frac{1}{2} \sum_{k} k v_{k} \text{ (double counting formula)}$$

$$\chi = |V| - \frac{n-2}{n} |E|$$

$$\bar{d} = \frac{\sum_{k} k v_{k}}{\sum_{k} v_{k}} = \frac{2|E|}{|V|} = \frac{2n}{n-2} \left(1 - \frac{\chi}{|V|}\right) \xrightarrow{|V| \to \infty} \frac{2n}{n-2}$$

to be compared with

$$\chi = \sum_{p \in V} K_c(p) = |V| \left(1 - d \frac{n-2}{2n}\right)$$

# Exceptional tilings of $\mathbb{T}^2$

The torus possesses regular tilings by triangles, quads and hexagons. An exceptional vertex is a vertex with different degree from the expected one (e.g. d = 6 for triangles).

Gauss-Bonnet implies : average degree remains constant (e.g. d = 6).

Theorem (Izmestiev, Kusner, Rote, Springborn & Sullivan)

There are no triangle tilings of the torus with two exceptional degrees (5,7).

Proof: uses geometric with equilateral (hence flat metric),  $\alpha = \pi d/3$ .

# Parallel transport and holonomy

If  $\gamma$  is a *smooth* curve on *S* (no curved vertex), then one can define parallel transport of vector *X* along  $\gamma$  (using unfolding)

If  $\gamma$  is a loop, the transported vector  $\tau_{\gamma}(X)$  is a rotated image of X, the same for any X. If  $\gamma$  is a loop around a vertex p, the rotation angle is exactly  $\alpha(p)$ .

The group generated by the rotations  $\tau_\gamma$  for all loops  $\gamma$  is the holonomy group.

#### Lemma

For the equilateral metric with identical faces (with n faces, n = 3, 4, 6), the holonomy is a subgroup of  $\mathbb{Z}_6$ ,  $\mathbb{Z}_4$  or  $\mathbb{Z}_3$ .

Proof.

- X does not change within a face
- ► the angle (X, edge of exit) changes from (X, edge of entry) by an integral multiple of <sup>π</sup>/<sub>3</sub> or <sup>π</sup>/<sub>2</sub> or <sup>2π</sup>/<sub>3</sub>

### Shortest geodesic

Let  $\mathbb{T}^2$  be endowed with the triangle equilateral metric and two exceptional vertices  $p_{\pm}$  of degrees 5 and 7 (angle defects  $\pm \frac{\pi}{6}$ . Consider  $\gamma$  a shortest geodesic (non nec. unique).

- $\gamma$  does not meet  $p_+$
- ▶ since  $\gamma$  does not meet  $p_+$ , it can be translated until it meets  $p_-$ , with side angles  $\pi + \phi, \pi + \psi, \phi + \psi = \pi/6$  ( $\phi, \psi \ge 0$ )
- ▶ if  $\phi$  or  $\psi$  are non zero, the holonomy is not in  $\mathbb{Z}_6$ ,  $\mathbb{Z}_4$  or  $\mathbb{Z}_3$
- otherwise suppose φ = 0, then one side of γ is flat: we can translate γ to γ' meeting again p<sub>−</sub> (and p<sub>−</sub> only)
- ▶  $\gamma, \gamma'$  delimit a flat cylinder, bounded by curves with positive angles  $\theta_1, \theta_2$ , such that  $\theta_1 + \theta_2 = \pi/6$ , with  $p_-$  on both boundaries
- ▶ let  $\delta$  link  $p_-$  to itself along the cylinder. Then a small perturbation of  $\delta$  will have  $\theta_1$  holonomy.